

A theoretical approach for node-ID negotiation in CANopen networks

Alexander Philipp, emotas embedded communication GmbH

The increasing complexity and size of modern CANopen and CANopen FD networks creates new challenges for system designers and device developers how to decide which CANopen node-ID should be used for their CANopen devices. This becomes more difficult in systems which are highly dynamic or having multiple devices of the same type in the network, like cascaded battery systems. Sometimes it is doesn't even matter for the system which node-ID a device has. This paper discusses a theoretical approach how devices in a CANopen or CANopen FD network could negotiate their node-ID by themselves.

Introduction

The assignment of a node-ID to a CANopen device is an essential requirement. Without a node-ID a CANopen device is unable to communicate over services defined in CiA 301 and CiA 1301.

Without a valid node-ID in the range of 1 to 127, a CANopen device would never leave the NMT sub-state reset communication. Since the node-ID has to be unique in the network, there has to be a way to configure the value for node-ID. Neither CiA 301 nor CiA 1301 specify a way to configure this value. There are various solutions in the market how this can be achieved.

Some manufactures use some hardware solution on their devices, like dip- or rotaryswitches or encode plugs. Others use some kind of proprietary software configuration. All of this solutions need more or less knowledge about the system where the devices are integrated.

Another way to configure the node-ID is specified in the document CiA 305, Layer Set t ing Services. This specificat ion describes services and protocols for identifying CANopen devices and assigning node-IDs, which can be used by a so called LSS Master or some configuration tool.

The CANopen Application profile for building door control, CiA 416, also specifies a procedure for claiming node-IDs by the devices themself. But besides that this procedure is slow, it is also patented.

The following theoretical approach for negotiation of node-IDs introduced by the engineers of the emotas embedded communication GmbH should be a starting point how to overcome most of the disadvantages of the other solutions.

Use cases, a short description

Masterless systems

Some systems are designed very flexible. That means that a system could be build with different combinations of devices, exactly meeting the needed requirements. Sometimes it could be enough that only two devices are need, sometimes much more and sometimes there are multiple of the same device type in a network.

An example of such kind of a systems are heating systems based on CANopen. They can include multiple sensors and controllers depending location and size where the system is installed.

When creating the devices for such a system, it is not clear which kind of device could be unique in the system, so it can operate as CANopen Manager. The CANopen Manager is the CANopen Master, which is responsible as NMT master, plus additional functionality for example LSS Master.

One way to assign node-IDs without a LSS Master is when integrating the system. But this leads to static system, without a real plug and play possibility. And in case some parts are added, or have to be replaced, a n

unused node-ID has to be assigned to the replaced part, which might be impossible for service technician.

Another example are battery clusters, which are capable of plug and play. The only devices in this network are the batteries themselves. There are physically all the same and have the same software. In such a system the node-ID is irrelevant to the system, but is required for CANopen communication.

Systems with a master

In generic CANopen systems with a master, there are also possibilities that the actual distribution of the node-IDs is not relevant for the functionality of the system.

Examples are the CANopen Application profile for energy management systems, CiA 454, or the CANopen Application profile for special-purpose car add-on devices, CiA 447.

These specifications use the LSS Fastscan, defined in CiA 305 to detect devices and assign node-IDs to them.

With LSS Fastscan devices can be only detected one after another, and in the worst case of completely unknown devices it will take 128 messages to verify the every single bit of the CANopen LSS address. With response-timeout of 10ms, it would take 1.28s to detect one single device.

New approach

The idea behind this approach is that, in systems which do not depend on the node-ID distribution, the devices negotiate their node-ID themselves.

The idea is not new, for example in the document J1939™-81 of the SAE®, there is a description of a so-called address claim procedure. In this procedure the devices in a CAN-based network negotiate the Addresses, depending on values of the NAME field.

The problem here is that this procedure uses the 8-byte Data Field of the CAN frame. This leads to the fact, that in case two or more devices are claiming the same Address at the same time could lead to collisions on the CAN bus.

In our approach we are trying to avoid that by not using the Data Field at all, all the

information exchanged between devices are encoded in the Identifier Field.

Requirements

The main requirements for this procedure is that the software of the devices is able to send and receive CAN Data Frames in Classical Extended Frame Format, CEFF. Because most modern CAN FD controllers support sending and receiving this kind of Data Frames, this approach works in CANopen FD as well.

That the software can distinguish between Extended and Basic Frame Format is important too.

Usage of the CAN Identifier

The CAN Identifier of an Extended Data Frame has 29 bits available. For our approach we only need 13 bits. The least significant 13 bits are used for the node-ID negotiation.

The 13 bits are divided into two fields, the least significant 8 bits are used as Data Field, and the most significant 5 bits as Data Code. The following table illustrates the usage of the CAN ID.

Table 1: 29 bit CAN ID usage

| Bit 28-13 | Bit 12-8 | Bit 7-0 |
|-----------|-----------|------------|
| reserved | Data Code | Data Field |

Using 13 bits of the 29 bits of an extended CAN ID means that 44% of the bits are required, but the following formula:

$$\frac{2^{13}}{2^{29}} = \frac{8.192}{536.870.912} = 0,00001525$$

shows that only 0,0015% of all possible extended CAN IDs are reserved for this service.

The Data Code defines the meaning of the content of the Data field. If the most significant bit is set to 1, the Data Field contains parts of the CANopen object 0x1018, Identity object. For this procedure to work, all 4 sub-indices are required.

The following table shows which part of the Identity object is transmitted with which Data Code.

Table 2: Data Code 0x1018 correlation

| Data Code | 0x1018 Data |
|-----------|-------------------------|
| b10000 | sub-index 1 - bit 0-7 |
| b10001 | sub-index 1 - bit 8-15 |
| b10010 | sub-index 1 - bit 16-23 |
| b10011 | sub-index 1 - bit 24-31 |
| b10100 | sub-index 2 - bit 0-7 |
| b10101 | sub-index 2 - bit 8-15 |
| b10110 | sub-index 2 - bit 16-23 |
| b10111 | sub-index 2 - bit 24-31 |
| b11000 | sub-index 3 - bit 0-7 |
| b11001 | sub-index 3 - bit 8-15 |
| b11010 | sub-index 3 - bit 16-23 |
| b11011 | sub-index 3 - bit 24-31 |
| b11100 | sub-index 4 - bit 0-7 |
| b11101 | sub-index 4 - bit 8-15 |
| b11110 | sub-index 4 - bit 16-23 |
| b11111 | sub-index 4 - bit 24-31 |

Due the binary coding of the Data Code, it is very easy to implement the algorithm for accessing the value of the Identity object. Bit 3 and 4 are the sub-index number - 1, and bit 0 and 1 are the byte number of this sub-index.

The following code fragment shows a simple implementation to create the corresponding CAN ID.

```

idx = byteNr / 4;
canId = identity[idx];
idx <<= 2;
idx += (byteNr % 4);
idx <<= 8u;

canId >>= ((byteNr % 4) * 8u);
canId &= 0xff;

canId |= idx;
canId |= 0x1000u1;

```

Two Data Code are for the flow control of the process and are listed in the following table.

Table 2: Data Code 0x1018 correlation

| Data Code | Name |
|-----------|---------------|
| b00001 | ReqUsedNodeId |
| b00010 | ActUsedNodeId |

The Data Code ReqUsedNodeId is used to start a new process, whereby the Data Field contains 0.

The resulting message with extended CAN ID 0x100 will trigger all devices with a valid node-ID to respond with ActUsedNodeId and their own node-ID in the Data Field.

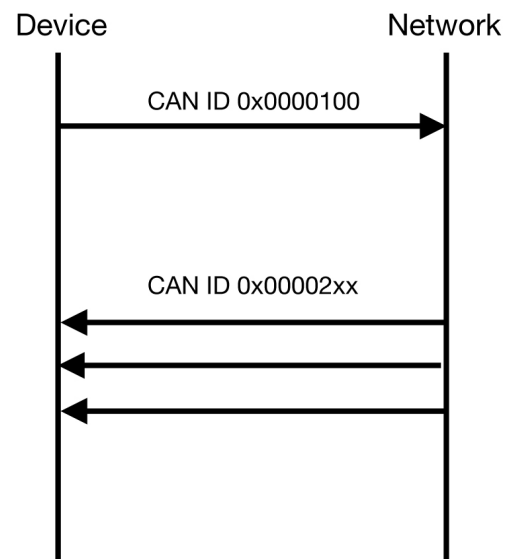


Figure 1: Initializing negotiation

Figure 1 shows the basic concept of asking the network for the node-IDs which are already used in the network.

Negotiation process

The negotiation process is only performed by nodes with an invalid node-ID.

All nodes with a valid node-ID are ignoring the messages with the most significant bit in the Data Code field set to one.

If a device participating in the procedure, receives a negotiation message it compares the value according to Table 2 with its own equivalent. In case the received value is higher than its own value, the device continues the negotiation.

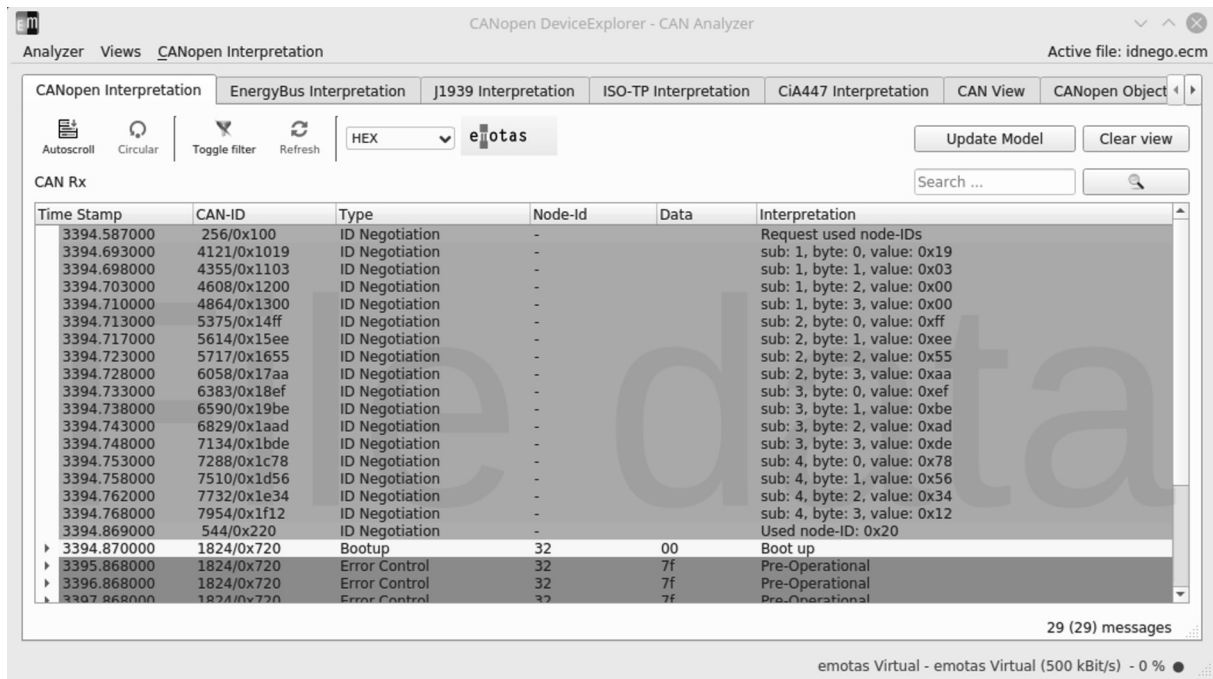


Figure 2: Complete negotiation cycle

In case the received value is smaller as its own equivalent, the negotiation is stopped on this device. And the device has to wait until the running negotiation is done. It can detect this by setting a timeout e.g. 100ms, and restarting this timeout every time it receives a negotiation message.

In case a device is able to transmit all 16 messages, which are needed to transmit all 128 bits of the Identity object, it will then announce the preferred node-ID to the network by using ActUsedNodeId Data Code message.

Figure 2 shows a complete negotiation cycle with timeout value of 100ms before and after the actual negotiation and an additional 5ms delay between every negotiation message.

Node-ID conflicts

In case a device successfully negotiated a node-ID, but another device indicates by the ActUsedNodeId message the very same node-ID, the receiving device sends out the same ActUsedNodeId and set its pending node-ID to the invalid node-ID, and enter the NMT state Reset Communication. Because of the other device receiving this message it loses the node-ID too.

Both devices can restart the negotiation process again. The losing device of this negotiation shall alter the preferred node-ID for the next cycle to a node-ID not yet present in the system.

Example

Table 4 shows a simulation example of two nodes negotiating their node-IDs. The difference in the 0x1018 parameter Revision Number causes the one node to stop the negotiation and restarts again after the first one is done.

Table 4: CAN Trace

| time | id | ide | l | d |
|----------|-------------|-----|---|------|
| 0.000000 | 256/0x100 | EXT | 0 | |
| 0.105000 | 4121/0x1019 | EXT | 0 | |
| 0.106000 | 4121/0x1019 | EXT | 0 | |
| 0.110000 | 4355/0x1103 | EXT | 0 | |
| 0.111000 | 4355/0x1103 | EXT | 0 | |
| 0.115000 | 4608/0x1200 | EXT | 0 | |
| 0.115000 | 4608/0x1200 | EXT | 0 | |
| 0.119000 | 4864/0x1300 | EXT | 0 | |
| 0.121000 | 4864/0x1300 | EXT | 0 | |
| 0.125000 | 5375/0x14ff | EXT | 0 | |
| 0.126000 | 5375/0x14ff | EXT | 0 | |
| 0.130000 | 5614/0x15ee | EXT | 0 | |
| 0.130000 | 5614/0x15ee | EXT | 0 | |
| 0.135000 | 5717/0x1655 | EXT | 0 | |
| 0.136000 | 5717/0x1655 | EXT | 0 | |
| 0.140000 | 6058/0x17aa | EXT | 0 | |
| 0.141000 | 6058/0x17aa | EXT | 0 | |
| 0.145000 | 6383/0x18ef | EXT | 0 | |
| 0.146000 | 6383/0x18ef | EXT | 0 | |
| 0.149000 | 6590/0x19be | EXT | 0 | |
| 0.155000 | 6829/0x1aad | EXT | 0 | |
| 0.161000 | 7134/0x1bde | EXT | 0 | |
| 0.165000 | 7201/0x1c21 | EXT | 0 | |
| 0.170000 | 7491/0x1d43 | EXT | 0 | |
| 0.174000 | 7781/0x1e65 | EXT | 0 | |
| 0.181000 | 8071/0x1f87 | EXT | 0 | |
| 0.279000 | 576/0x240 | EXT | 0 | |
| 0.280000 | 1856/0x740 | _ | 1 | 0x00 |

| | | | | |
|----------|-------------|-----|---|------|
| 0.282000 | 256/0x100 | EXT | 0 | |
| 0.283000 | 576/0x240 | EXT | 0 | |
| 0.387000 | 4121/0x1019 | EXT | 0 | |
| 0.392000 | 4355/0x1103 | EXT | 0 | |
| 0.397000 | 4608/0x1200 | EXT | 0 | |
| 0.402000 | 4864/0x1300 | EXT | 0 | |
| 0.406000 | 5375/0x14ff | EXT | 0 | |
| 0.412000 | 5614/0x15ee | EXT | 0 | |
| 0.417000 | 5717/0x1655 | EXT | 0 | |
| 0.422000 | 6058/0x17aa | EXT | 0 | |
| 0.426000 | 6383/0x18ef | EXT | 0 | |
| 0.432000 | 6590/0x19be | EXT | 0 | |
| 0.437000 | 6829/0x1aad | EXT | 0 | |
| 0.442000 | 7134/0x1bde | EXT | 0 | |
| 0.447000 | 7423/0x1cff | EXT | 0 | |
| 0.451000 | 7491/0x1d43 | EXT | 0 | |
| 0.456000 | 7781/0x1e65 | EXT | 0 | |
| 0.462000 | 8071/0x1f87 | EXT | 0 | |
| 0.561000 | 577/0x241 | EXT | 0 | |
| 0.562000 | 1857/0x741 | _ | 1 | 0x00 |
| 1.280000 | 1856/0x740 | _ | 1 | 0x7f |
| 1.562000 | 1857/0x741 | _ | 1 | 0x7f |

Summary

The main goal is still to find a reliable, robust, and fast solution to assign node-IDs in plug and play systems to extend the network during its lifecycle. This theoretical approach should be only a basis for discussion of how to achieve this.

With the discussed process it is possible to assign node-IDs to devices without the need of a dedicated LSS Master. The usage of extended CAN IDs makes it possible that only 16 messages are needed, and that the process doesn't create CAN-BUS collisions.

Alexander Philipp
emotas embedded communication GmbH
Fritz-Haber-Str. 9
DE-06217 Merseburg
www.emotas.de

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